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The Tao of Teiid

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Sept. 4, 2009



What is Teiid?

- Teiid is an open source solution for scalable information integration through a relational abstraction.
- Teiid focuses on:
 - Real-time integration performance
 - Feature-full integration via SQL/Procedures/XQuery
 - Providing JDBC access
- Teiid enables:
 - Data Services / SOA
 - Legacy / JPA integration



Overview

- Background
- Architecture
- Internals
- Wrap-up
- Q & A



Where did Teild come from?

- Project lineage is from MetaMatrix starting in ~1999.
 - Teiid http://www.jboss.org/teiid
 - Teiid Designer http://www.jboss.org/teiiddesigner
 - DNA http://www.jboss.org/dna/
- MetaMatrix was the leader in Enterprise Information Integration (EII) – hence Teiid.
- Red Hat acquired MetaMatrix in 2007.
- Last major MetaMatrix product release, 5.5.3 10/2008



Project Status

- Open source 2/2009 heavily refactored from 5.5 line
- 6.0 Initial release 3/2009
- 6.1 Teiid / Teiid Designer release 6/2009
- 6.2 Coming Soon! Embedded/server deployments, Designer-less usage, AdminShell, and much more.
- Anticipate a platform release combining Teiid and other JBoss technologies next year.

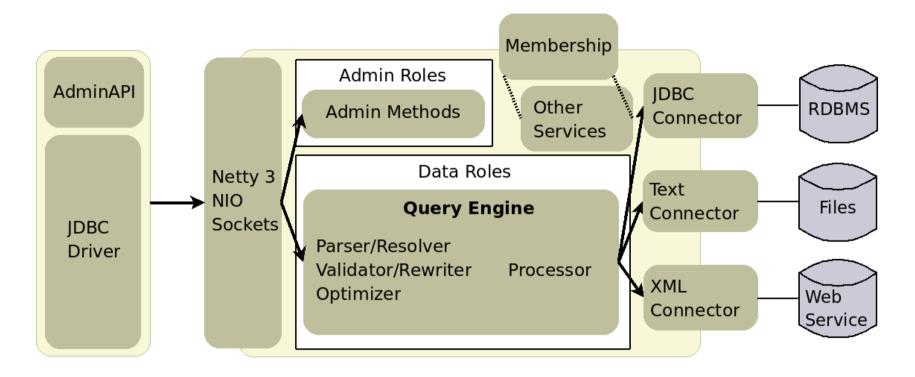








Architecture



- SEDA Connector bindings, socket transport, query engine, admin methods all have queues/thread pools
 - Each connector binding operates independently
- Other services include JBoss Transactions JTA, BufferManager, sessioning, etc.

Connector API

- Simplified object form of JDBC with concepts of JCA.
 - Pooling, caching, some security handled by the runtime.
 - Queries are resolved objects not just a string.
- Extended metadata (ConnectorCapabilities) directs the optimizer source query formation.
- In addition to out of the box offerings, our JDBC Connector is easily extended.
- Can be thought of as a JDBC toolkit.



Other Extension Points

- Logging (Log4j), specific contexts for audit and commands
- MembershipDomains handle authentication/group assignment. Provide File and LDAP by default.
- User defined functions Implementation method in Java, currently only defined through Designer.
- Scripting through AdminShell







Teiid Internals

- Integration Features
- Planning
- Processing
- Transactions



Integration Features

- Access Patterns criteria requirements on pushdown queries
- Pushdown decompose user query into source queries
 - Remove unused select clause items
 - Decompose aggregates over joins/unions
- Dependent Joins (can use hints) feed equi-join values from one side of the join to the other
- Optional Join (can use hints) removes an unused join child
- Multi-source connector bindings allows for multiple homogeneous schemas to be used through the same model.
- Copy Criteria uses criteria transitivity to minimize join tuples.

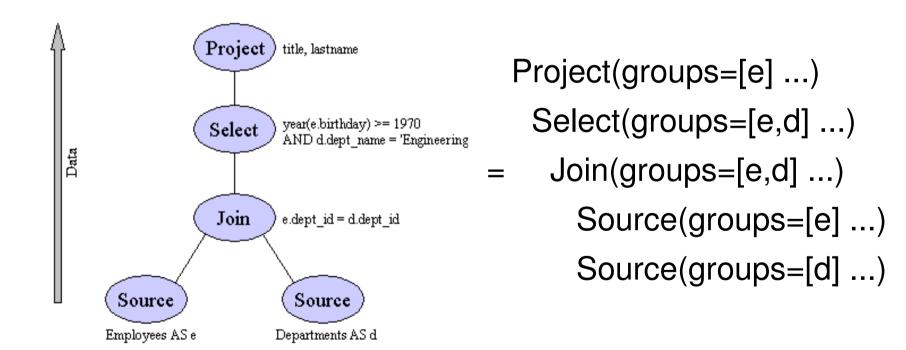


Planning

- Distinct phases: parse, resolve, validation, rewrite, optimization, process plan creation.
- Rewrite canonicalizes and simplifies.
- The optimization phase follows with rules/hints/costing
 - Procedures/XQuery not formally optimized
 - Non-federated optimization is similar to mature RDBMS
- Optimizer plan structure is a flexible tree distinct from the command form and processing plans.
- Planning is typically quick and deterministic prepared plans are recommended

Visualizing a Plan

select e.title, e.lastname from Employees as e JOIN
Departments as d ON e.dept_id = d.dept_id where
year(e.birthday) >= 1970 and d.dept_name = 'Engineering'





Understanding Planning

 Initial canonical plans follow the logical SQL processing flow:

from/where/group by/having/select/order by/limit/into

- Each node corresponds to a logical SQL operation
- Canonical relational plans not performant for federated queries – optimization is necessary
- Processing plans and intermediate plans can be shown in the log/obtained by the client.

```
select * from ... option debug - With DETAIL logging
```



Plan Rules

- Initial sequence driven by query form some rules trigger others
- Move/create/delete/modify nodes toward more optimal form
 - RemoveVirtual Removes inline views or nested transformations
 - RaiseAccess Ensures access nodes are raised meaning more will be executed by the connector
 - PushSelectCriteria Moves criteria toward tuple origin
 - CollapseSource Takes plan nodes below an Access node and creates a query (not the final query sent to the source, which will get translated by the connector)
 - RulePlanSorts Combines sort processing operations ... many others ...
- Many rules correspond directly to federated optimizations –
 CopyCriteria, AggregratePushdown, RemoveOptionalJoins, etc.



Example Rule Application

```
SetOperation(groups=[], props={USE ALL=true, SET OPERATION=UNION})
 Project(groups=[BQT1.SmallA], props={PROJECT_COLS=[IntKey]})
  Access(groups=[BQT1.SmallA], props={MODEL ID=Model(BQT1)})
   Source(groups=[BQT1.SmallA], props={NESTED COMMAND=null})
 Project(groups=[BQT1.SmallA AS SmallA 1], props={PROJECT COLS=[SmallA 1.IntNum]})
  Access(groups=[BQT1.SmallA AS SmallA 1], props={MODEL ID=Model(BQT1)})
   Source(groups=[BQT1.SmallA AS SmallA 1], props={NESTED COMMAND=null})
EXECUTING RaiseAccess
AFTFR.
Access(groups=[], props={MODEL ID=Model(BQT1)})
 SetOperation(groups=[], props={USE ALL=true, SET OPERATION=UNION})
  Project(groups=[BQT1.SmallA], props={PROJECT_COLS=[IntKey]})
   Source(groups=[BQT1.SmallA], props={NESTED COMMAND=null})
  Project(groups=[BQT1.SmallA AS SmallA 1], props={PROJECT COLS=[SmallA 1.IntNum]})
   Source(groups=[BQT1.SmallA AS SmallA__1], props={NESTED_COMMAND=null})
```



Join Planning

- The most complicated parts of the optimizer
- It is not exhaustive, but does consider ordering (left linear), satisfying access patterns, and algorithm ([Partitioned] Merge / Nested Loop)
- Ordering/algorithm is only important for federated joins.
 Once a join is pushed, it's declarative to the source
- Merge joins have dependent variants, which can have large impact on performance – especially an unnecessary dependent join (see makenotdep)



Use of Costing

- Specified as attributes at the table and column level will have a runtime interface soon
- Mostly based on cardinality with a simplistic cost model of execution
- Assign costs to different join ordering and implementations to pick the best one
- Using small, or inappropriate values, could lead to unexpected performance
- See plan info "Estimated Node Cardinality", "Estimated Independent/Dependent ...", etc. for values used in planning.

Processing

 A relational processing plan is composed of discrete operations organized as a tree – very similar to the optimizer form:

```
AccessNode – Source Query/Procedure

GroupingNode – Grouping operations and aggregate calculation

JoinNode – Joins the left and right tuple sources together

LimitNode – Honors limits and offset

ProjectNode – Converts tuples (select clause)

SelectNode – Applies selection (where clause) criteria

SortNode – Sorts incoming tuples
...
```

- Procedure plans are composed of instructions.
- Tuples are processed in batches. The BufferManager is set to a specific memory limit; excess batches are written to disk.
- Processing algorithms are sort based, variants chosen during planning and processing.



Example Process Plan

select * from System.DataTypeElements

```
ProjectNode(1) [dt.Name AS DataTypeName, c.NAME, ...]
JoinNode(2) [PARTITIONED SORT JOIN (SORT/SORT_DISTINCT)] [INNER JOIN]
\text{criteria=[c.PARENT_UUID=dt.UID]}

AccessNode(3) SELECT c.PARENT_UUID, ... FROM SystemPhysical.COLUMNS AS c
ProjectNode(4) [dt.NAME, dt.IS_BUILTIN AS IsStandard, ...]
JoinNode(5) [MERGE JOIN (SORT/SORT)] [LEFT OUTER JOIN]
\text{criteria=[dt.UUID=a.ANNOTATED_UUID]}

AccessNode(6) SELECT dt.UUID, dt.NAME, ... FROM SystemPhysical.DATATYPES AS dt
AccessNode(7) SELECT a.ANNOTATED_UUID, ... FROM SystemPhysical.ANNOTATIONS AS a
```

- Shows decomposition into 3 source queries.
- Also the optimizer has combined a distinct operation into JoinNode(2) loading of the right child.



Handling Load

- Memory Usage the BufferManager acts as a memory manager for batches (with passivation) to ensure that memory will not be exhausted.
- Non-blocking source queries rather than waiting for source query results processor thread detach from the plan and pick up a plan that has work.
- Time slicing plans produce batches for a time slice before re-queuing and allowing their thread to do other work (preemptive control only between batches)
- Caching ResultSets at the connector and user query level can be reused on a session or vdb basis



Transactions

- Three scopes
 - Global (through XAResource)
 - Local (autocommit = false)
 - Command (autocommit = true)
- All scopes are handled by JBoss Transactions JTA
- Command scope behavior is handled through txnAutoWrap={ON|OFF|OPTIMISTIC|PESSIMISTIC}
- Isolation level is set on a per connector basis.









Performance

- Raw (cpu-intensive) overhead is typically submillisecond per prepared user query.
- Integration performance check the processing plan.
 We'll usually have the best form.
- Consider using UDFs (Java) for reusable subroutines rather than stored procedures.
- Client result sets can be scroll insensitive and backed by the BufferManager.



Differences with traditional Java DBs

- Flexible planning architecture
- Geared to high-performance integration processing task specific queues and thread pools, advanced buffer management, batching, etc.
- Lack of DDL support
- Loose constraint handling
 - pk/fk, unique, and type constraints are in metadata, but are not enforced at runtime.
- Temp tables backed by BufferManager rather than a relational/indexed storage engine.



Future Releases

- We'll look even more like a database direct usage of DDL for metadata.
- More features around materialization, data locality, and caching.
- Continued integration with other JBoss projects.
- More design-time integration with Eclipse DTP http://www.eclipse.org/datatools/



QUESTIONS?

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